

APPLYING MULTI- LEVEL ANTCOLONY OPTIMIZATION TO RESTRICTED FOREST



Prashant Kumar

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University Department of COMPUTER SCIENCE, B.R.A. Bihar University, Muzaffarpur, India.

E-mail: prashantjha1983@gmail.com

ABSTRACT

Since the beginning of the last century, one of the most important concerns has been the problems arising from the planning of forest transport. Off-road activities are those that involve the transportation of wood from stump location to roadside landings or landings in centralized locations. The term "on-road activities" refers to the process of transporting timber to its final destinations using land-based vehicles.

Exact algorithm and approximation algorithm, both are different methods which have been used in the process of solving ftp. The core of approximation algorithms, commonly referred to as heuristics, is the process of ranking a large number of possible answers and choosing the most appropriate one. The fact that exact algorithms can generate optimal answers is the most important advantage associated with using them.

KEYWORDS: Multi- Level, Ant colony.

INTRODUCTION

Integer and combined integer fashions can model transportation issues better than continuous variable models. However, due to the apparently high computing complexity of IP and MIP models, their use is limited to the solution of small to medium-sized problems (Weintraub 1995; Olsson 2003). In addition, trends in geographic record structures (GIS) have made it easier to create and manipulate information covering substantial areas, making it easier to build large-scale challenges. Furthermore, many FTTPs do not have a formal mathematical description, so it is impossible to apply precise strategies to them (Murray 1998). Consider undertaking the construction of a road community in a densely wooded area in such a way that it allows easy access to some timber income while at the same time reducing the overall value of road production. This problem, known as the multiple goal entry to problem (mtap) using Dean (1997) and can only be solved using heuristic techniques analogous to Murray (1998).

Due to the fact that at the beginning of the last century, one of the most full-sized concerns has been the problem that arises from the planning of wild area transportation. In general, forest transportation planning problems (FTTP) can be divided into off-avenue and on-avenue degrees, each of which may be hypothetically dependent on each other (Heinemann 2001). Off-avenue haulers are those who transport lumber from stump locations to roadside landings or landings in centralized locations. The time term "on-road sports" refers to a land-based method of transporting timber to its final locations entirely using cars.

Exact algorithms and approximation algorithms, each of which are terrible methods, were used in the process of fixing ftp. The core of approximation algorithms, commonly known as heuristics, is the technique of ranking a large number of possible solutions and choosing the most suitable one. The truth is that specific algorithms can generate the most complete answers, with the maximum benefit associated with their use. But, they may be best applied to problems of a restricted scope. Alternatively, heuristic procedures, regardless of the fact that they may not yield excellent feasible answers, were successfully employed to overcome large scale problems and unique algorithms have been developed.

Due to the discrete nature of FTPP variables that include road construction, integer and combined integer fashions (IP and MIP) have attracted much interest in the last few years. This is because integer and mixed integer fashions can constitute transport perturbations better than non-stop variable models. However, due to the fairly high computing complexity of IP and MIP models, their use is constrained to the solution of small to medium length problems. Furthermore, many ftpp no longer have a proper mathematical description, consequently it is very impossible for them to use the correct techniques (Murray 1998). Recall the undertaking of building a road community in a heavily forested area in such a way that it allows clean access to positive timber income while at the same time minimizing the overall cost of road construction. This difficulty, which has been called the multiple round access problem (mtap) by Dean (1997) and according to Murray (1998), can be easily solved through the use of heuristic techniques. Similarly, some techniques have also been created that combine MIP with heuristic methods. While these methods are intended to capture the benefits of both methods, they provide partial solutions to the largest while increasing the efficiency of the exact algorithms. This answer results in a trade-off between great and efficiency, which is presented using heuristics (given by the actual algorithm). This usually happens for two reasons. To begin with, there may be no formal mathematical formulation capable of appropriately representing the difficulty of the task, which is clearly dependent on the forms of variables and goals that may be considered. Problems that arise in the 2D, real world are regularly of a scale that makes it impossible to deal with them effectively using the gear now available to find precise answers. Several algorithms that use heuristics were created to be a good way to measure ftpp in which the transport cost is either constant or variable. It has been performed so that it overcomes the constraints of unique methods (Chung and Sessions 2003).

Trouble reporting

Considering each constant and variable fee following the aspect rules became the object of this research challenge, which led to the development of a unique approach to problem solving, which is mainly based on the Eko metaheuristic. The mission that needs to be solved is to determine the set of routes that have the lowest common cost, from wood sales to generators, chosen as the very last destination, all within the wooded area road network. In the form of sediment yield considering environmental impacts. This unique subject matter, in conjunction with the vast majority of other transportation problems, is one that can be

defined as a network programming problem. A graph G is used to represent the road network system. These variables can be interpreted as representing distance, charge, or something different of the edge. As an end result, a network denoted through the letter n is created to symbolize the transportation planning problem. There are three specific values that pertain to each edge of this specific FTTP that is being taken into account. Those prices are variable price, fixed price and quantity of sediment production. The number of customers without delay each day affects the variable prices. These costs are fixed in nature. In the same way that constant costs are assumed, we assume that whenever roads are used, regardless of traffic, sediment builds up.

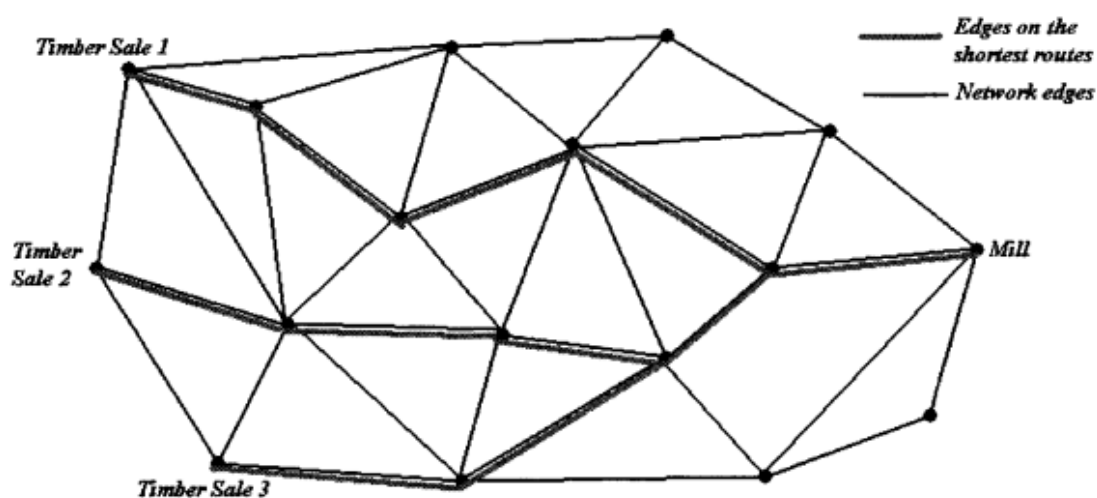


Figure 1. Consider the case of three timber sales and one mill site as an example of a transportation issue

Research Methodology

As established in the rest of the bankruptcy, a multiscale technique was successfully used for the parameter configuration problem.

Multilevel Planning

Typically, multilevel plans solve a large problem by using increasingly smaller issues harder and faster through a sequence of solution refinements.

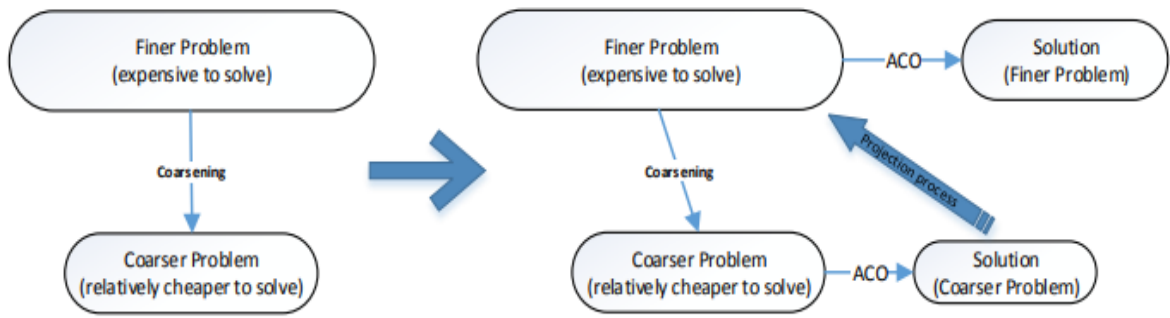


Fig. 2: Diagram showing a multilevel scheme

Multilevel Approach

Getting to the Coarse Level Problems

Since CFTPP's networks are weighted graphs, graph thickening strategies can be explicitly applied to supply the coarseness degree issues.

Built-in Echo

The underlying Echo set of rules (Algorithm 7), known in this look as "GuidedEcho", is designed for the solution of a fine-phase problem guided by means of interpolated solution components, which are calculated as Can be mapped from north. Serious level of trouble. Inside the implementation, GuidedACO is used for basic perturbation with or without guidance for all coarse level problems and interpolated answer components.

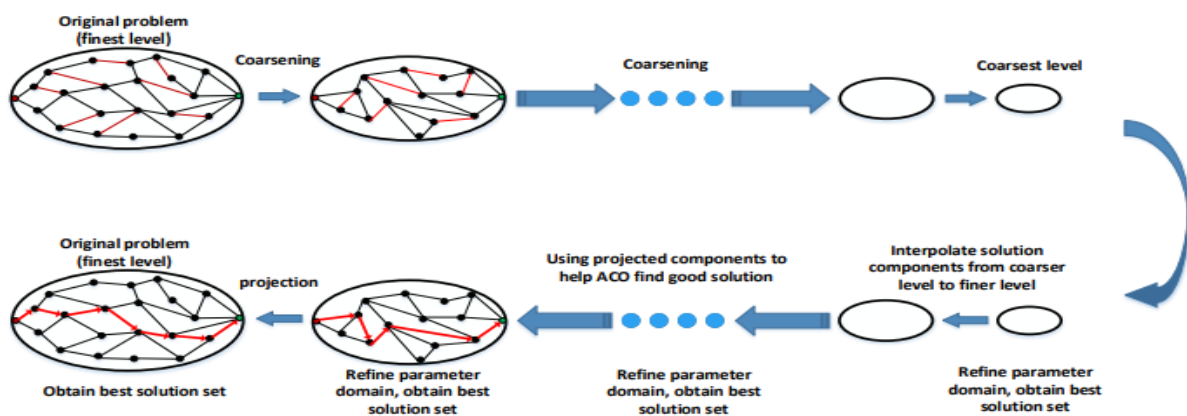


Figure 3: The top left corner

RESULTS

It picks up bankruptcy where bankruptcy five left off by means of continuing the investigation into the introduction and use of MOACO with the aim of resolving BofPP. In order for you to reduce the dependence of the aco algorithm on particular parameter settings, moaco is designed with earlier awesome algorithmic components than those found in current moaco algorithms. Even when Moaco was able to solve all test cases within the research, the amount of computing time that was critical to accomplish this varied greatly depending on the containment conditions and the timing of the test. Consequently, in this chapter, we layout a multilayer MOACO (mimoACO) so that it boosts MOACO's overall performance both in terms of the first rate of solutions it produces and the amount of time it takes to produce those products. In a way analogous to earlier works, mmoaco is first developed to solve problems at a coarse level from the initial perturbation to a coarse level by using rough heuristics, followed by Let's move on to solving the initial problem. Based on the findings of previous experiments on a wide range of these problems, it is quite economical to predict that MMOCO will demand to spend less time for high quality transactions than MOACO. In addition, Eko uses solutions located to problems at a coarse level if you want to help it discover possible answers to problems at a finer level, which further complements both the convergence speed and the answer. A description of this is more in depth and can be seen below.

Multilevel Ptero Echo

Rough precision graphs

For advanced multiscale schemes, it has been found from previous studies that milieu is important for problems at the micro- and coarse-scale levels in order to percentage some unusual habitats. This is essential so that the answers to issues at the coarser levels can be used efficiently to facilitate the process of finding finer level solutions that are of extremely high quality. For this reason, the coarsening method (set of rule 12) is performed to supply coarse degree networks that may be finer replicas of graphs in terms of general feature weights (sums of fees and sediments on each facet), but size in small. [C]orsencing methods produce networks that are smaller in size but are accurate replicas of finer graphs (less diversity of nodes and edges). This is achieved by adding the weights of the mixed edges to the aggregated nodes and preserving the entire weight of the collapsed edges. This is a departure from the method used in previous chapters, in which it was assumed that nodes had no weights associated with them.

Algorithm 12 Coarsening_Procedure(G)

$QE \leftarrow$ matching edges in graph G

foreach $edge(u,v) \in QE$ **do**

$u.v \leftarrow$ aggregate u and v

$weight(u.v) \leftarrow weight(u) + weight(v) + weight(u,v)$

if u, v are both adjacent to a node k **then**

$weight(u.v,k) := weight(u,k) + weight(v,k)$

end

end

So that it can find the edges that fit, we have perfected the Heavy Part Matching (HEM) algorithm. This set of rules calculates the maximum number of matches that contain edges with the highest weight. The nodes of a community are traversed alternately, one after the other, in order.

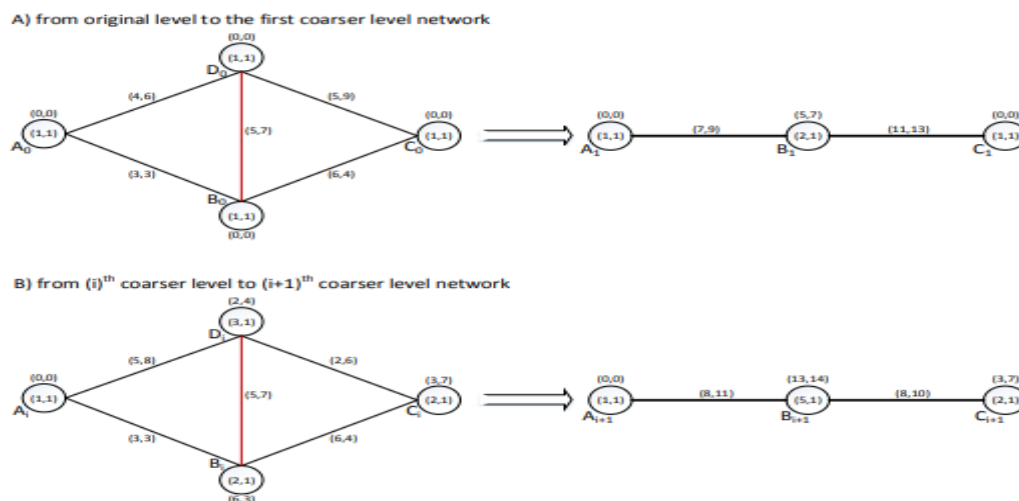
In the context of BotPP, the associated costs and sediment values are summed collectively to determine the weights that can be assigned to each edge. There are distinct advantages to using hempp. To start with, given that botpp's 2 number one goals are normal values and lack of sediment, it is very likely that answers will involve the use of thin weighted edges. HEM detects matching edges that have high weights at a more granular level of community. After collapsing the matching edges, the coarse phase graph that is produced will consist of low-weight edges instead of the original edges. Hem guarantees that the matching process is determined by traveling from one node to the following in a predetermined order and decides on the rims with the best overall weight. This removes the element of chance from the system's coarsening of graphs, which ends up in answers that are more reliable averages.

Aggregation and Importance Factors

The system starts by assigning facet charges and sediment values based on the initial perturbation. For each node the price and sediment price are both set to zero as the starting price. Importance and aggregation elements are one set for each, due to the initial cost. An aggregation factor of 1 means that no aggregation was performed. During the full coarsening process, the importance element will always be 1, and the aggregation element of the new node can be calculated as the sum of the aggregation elements of the merged nodes. This system is iterated over successive layers which can be coarsened with new values to be derived from the aggregation.

Answer Mapping Method

After coarse software for the initial network and gradually with an extended range of coarseness after a certain network is built up, an underlying set of rules (to be elaborated upon in later sections) begins to be used. The network with the highest degree of coarseness is used to search for answers.



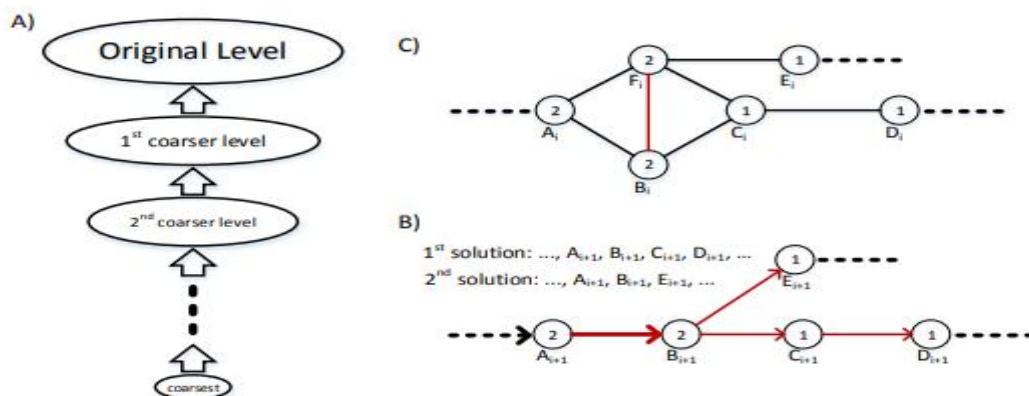
(4)

A charge (its amount is omitted in parentheses on each side) and the precipitate are both blanketed as field properties (the correct charge in parentheses on each side). A node's properties include its cost (the cost to the left of each node's parentheses), its sediment (the cost to the right of each node's parentheses), its aggregation factor (the charge to the left of each node's parentheses), and its properties. Are included.point of significance (exact value in parentheses at each node).

degree network. The answer seen for a low level problem is useful for the echo approach, it is used to solve a finer level problem because the answer is being used within the search system. This support is provided by means of a mapping technique, which interprets coarse level answer components into interpolated finer degree components (in a way this is similar to the technique used in bankruptcy 4), for which increasing solutions Additional considerations are provided for

At some rough level, the number of instances of the different types to be preserved for each node inside the answer set is enumerated and recorded. These statistics are then used to

establish significance, which is an important part of the method. The importance of nodes that are not part of the solution set is determined by



Set 4.2: A diagram showing the process of solution mapping as well as the computation of the critical elements, the $(i + 1)$ th coarse level community nodes B_{i+1} and A_{i+1} are of importance two, While the key point of the alternate nodes is best 1. This is because those nodes were protected in two different answers. The interpolated nodes of the i th level network employ the same number of importance factors as the original nodes. All components are similar to at least one. Mapping methods break down the answer into components at a coarse level in a finer phase of the problem. Interpolated problem additives are, in essence, finer stage (second thickest degree) problem additives, which can be collapsed or copied without delay to shape the acquired coarse stage solutions. This is because the interpolated perturbation additives are derived from the acquired coarse degree answers. During the technique of mapping, the range of examples in which each node in the solution set is preserved is matched and used to outline the important aspect. It is implemented so that the significance factor is as it should be decided.

Interpolated additives that are attached to nodes that contain elements of higher importance may be given higher concerns to be used as answer components. This is due to the fact that it is very expected that it can help to rapidly convert the built-in aco to the very first-class answer set.

Solution generation

Build a bachelor solution

The Pareto frontier is divided into several classes using the transition probability method, which uses a parameter specified using a value from zero to one in an incremental fashion. The relative importance of dreams can be adjusted using manipulation of the values of α . If the value of α is both zero or 1, then the algorithm has a greater tendency to produce solutions that are directed towards each objective function (i.e., the value or sediment). However, if the value of α is a piece with a price inside the $[0, 1]$ type, more emphasis is placed on answers that involve trade-offs. Moaco investigates the capacity solution by propagation of Pareto spheres. The answers set for each unique field are combined to arrive at the final result.

Conclusion

In the years to come, I hope to continue the lines of inquiry I am working on now. In particular, I'd love to continue the study I started in multilevel planning, as well as the studies I started on hierarchical computing, and I'd also like to expand on the multilevel structure that I've Made for additional domains and customizations. Technique. The following elements provide a brief rationalization of my ideals, along with several courses of action I plan to take.

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